

# **Organisatorisches (1)**

- **Lehrstuhl „Echtzeitsysteme und Kommunikation“**

- <http://ivs.cs.uni-magdeburg.de/EuK>
  - [nett@ivs.cs.uni-magdeburg.de](mailto:nett@ivs.cs.uni-magdeburg.de)
  - Sekretariat Frau Duckstein, 29/405, Tel. 67-18345

- **Web-Adresse**

<http://ivs.cs.uni-magdeburg.de/EuK/lehre/lehrveranstaltungen/>

- Folien der Vorlesung (in englisch)
  - praktische Übungsaufgaben
  - Mitteilungen (z.B. Bekanntgabe der Übungsgruppeneinteilung)
  - Literaturhinweise

- **Übungen**

- Übungsleiter: Herr Lindhorst (E-mail: [lindhors@ivs.cs.ovgu.de](mailto:lindhors@ivs.cs.ovgu.de))

Herr Dornemann, Herr Penzlin

- Wöchentlich ab 15. KW (13. April)

# Organisatorisches (2)

- **Inhalte der Übungen:**

- Vermittlung praktischer Kenntnisse in der Anwendung

Programmieren von TCP-Sockets. Während des Semesters wird in der Übung ein Webserver programmiert, der weiterhin einen Chat-Service anbietet. Dazu werden in regelmäßigen Abständen Teilaufgaben zu bearbeiten sein. Die Programmierung erfolgt unter Linux in C. Grundlegende C-Kenntnisse werden vorausgesetzt, bzw. die Bereitschaft, sich solche anzueignen.

- Wenn erforderlich, gewünscht ....: Aufarbeitung des Vorlesungsstoffs

Die Übungen finden in unserem Labor statt (Raum 425).

- **Kriterien für erfolgreiche Teilnahme:**

- schriftliche Prüfung (Scheinerwerb: Note <= 4.0)

- Zulassungsvoraussetzungen:

- erfolgreiche Umsetzung der praktischen Aufgaben, sowie Vorstellung und Erklärung der Implementierung

# **Introduction (1)**

## **Computer network:**

An *interconnected* collection of *autonomous* computers

## **Interconnected computers:**

Computers are said to be interconnected if they are able to exchange information

## **Examples for the physical connection medium:**

copper wire, fiber optics (wired)

radio communication (micro waves), infrared (wireless)

## **Autonomous computers:**

No computer can forcibly start, stop, or control computations (actions) on another one.

## **Distributed system:**

The existence of multiple autonomous computers is transparent (not visible) to the user,i.e. no explicit path addressing (e.g. URL), no explicit moving of objects (files) etc.

## **Principle vehicle for discussing computer networks and their protocols:**

The Internet

# What's the Internet: “component's” view (1)



PC (host)



server



wireless laptop



cellular handheld



access points



wired links



router

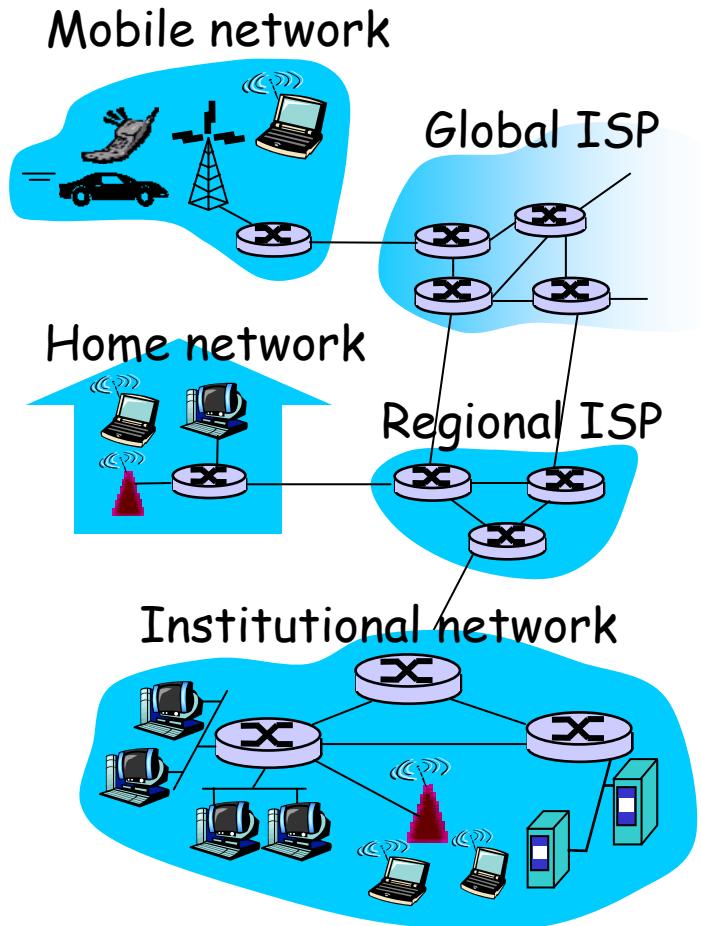
- millions of connected computing devices:

*hosts = end systems running network applications*

- *communication links*

- fiber, copper, radio, satellite
- transmission rate *bandwidth*

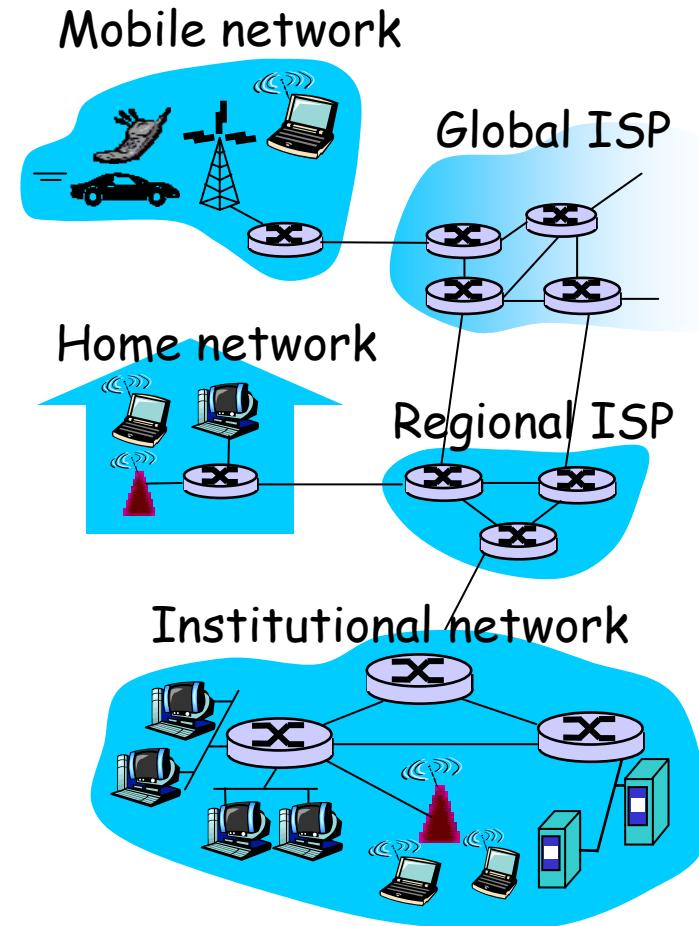
- *routers*: forward packets (chunks of data)



## What's the Internet: “component’s” view (2)

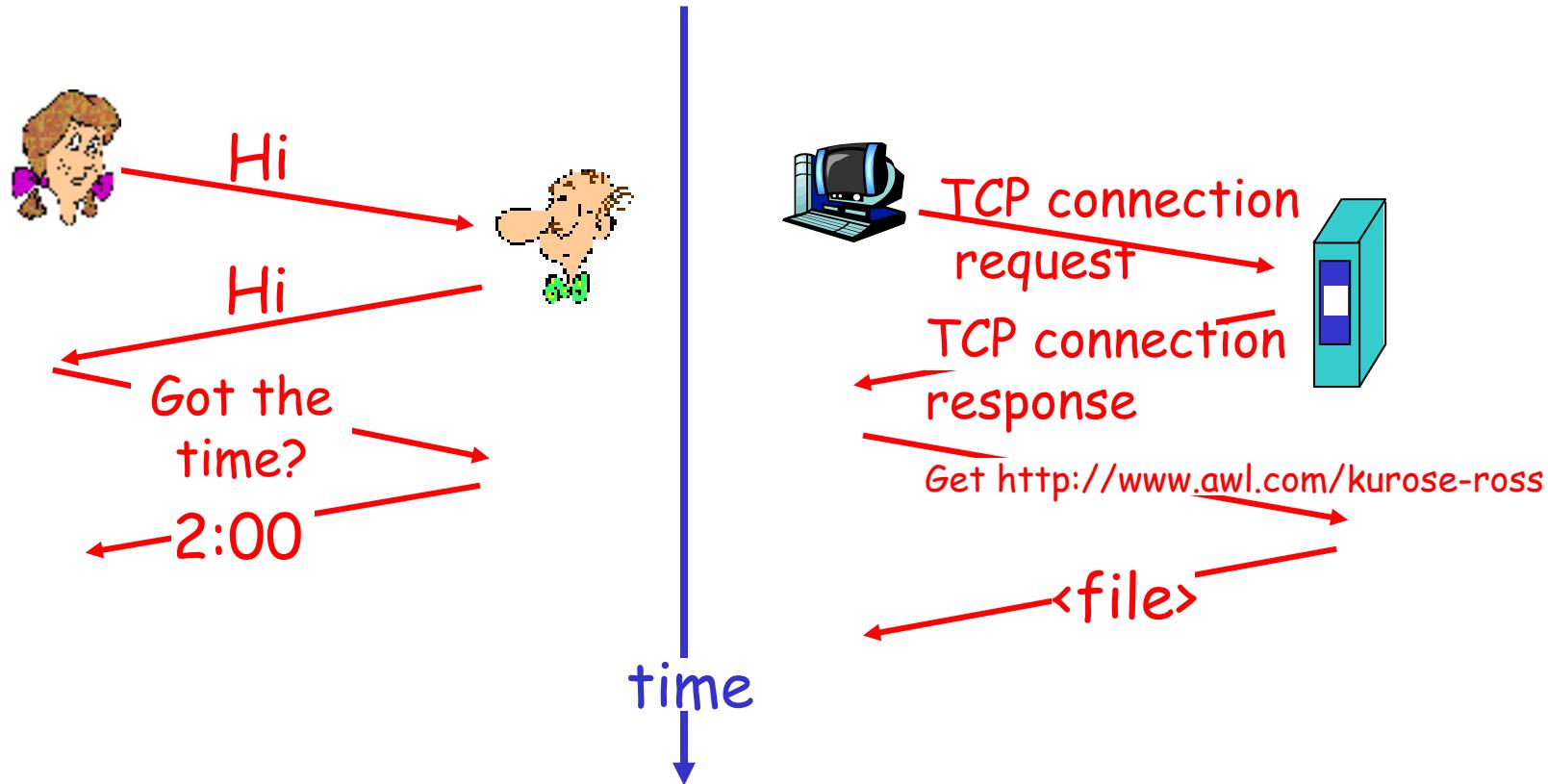
- **protocols** control sending, receiving of msgs
  - ❖ e.g., TCP, IP, HTTP, Ethernet

- Internet standards
  - ❖ RFC: Request for comments
  - ❖ IETF: Internet Engineering Task Force



# What's a protocol?

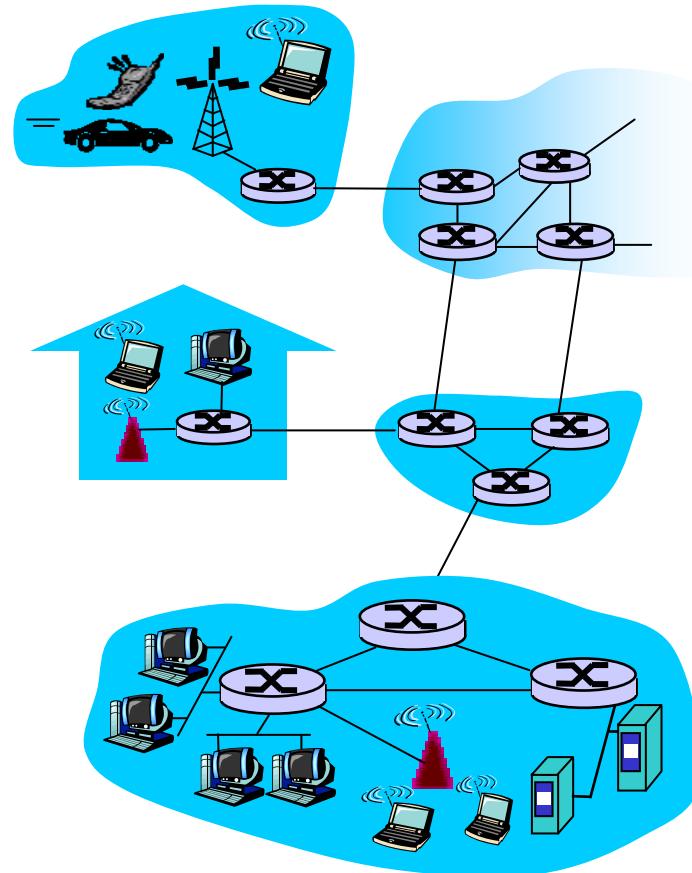
a human protocol and a computer network protocol:



*protocols define format, order of msgs sent and received among network entities, and actions taken on msg transmission, receipt*

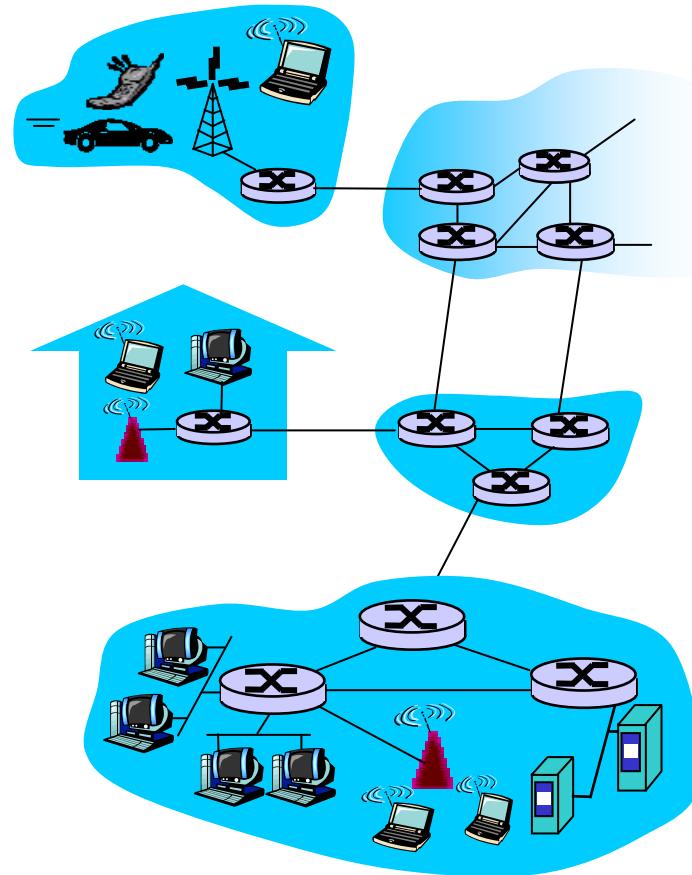
# What's the Internet: a service view

- communication *infrastructure* enables distributed applications:
  - Web, VoIP, email, games, e-commerce, file sharing, Skype
- communication services provided to apps:
  - reliable data delivery from source to destination
  - “best effort” (unreliable) data delivery



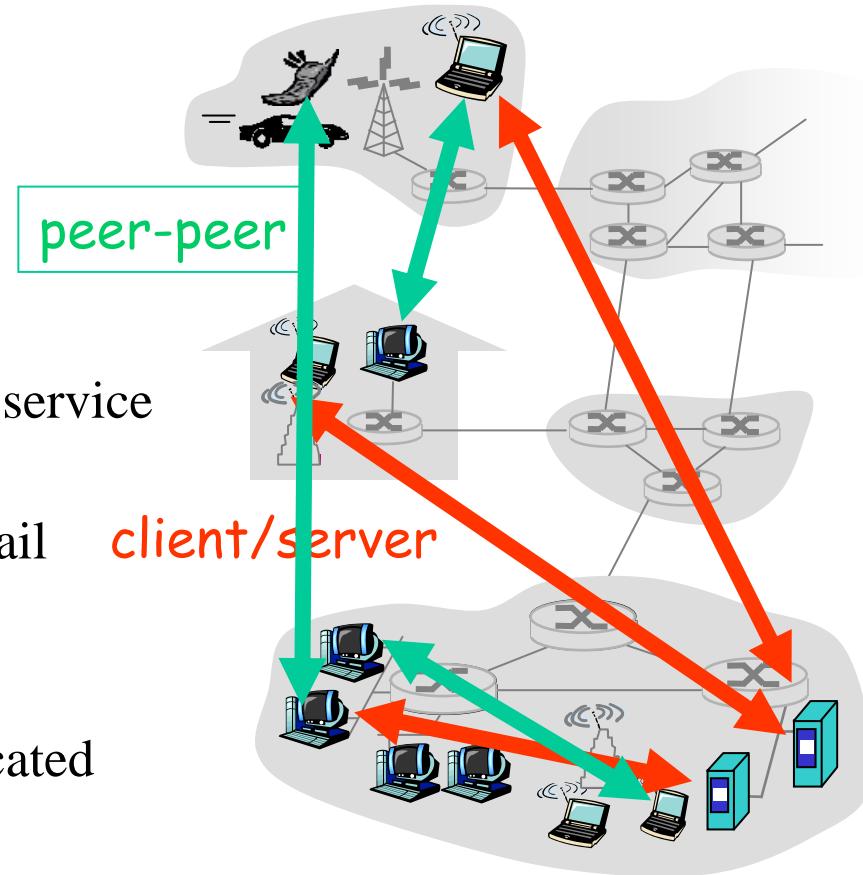
# A closer look at network structure

- **network edge:**  
applications running on hosts
- **access networks:**  
wired or wireless communication  
links that connect hosts to their edge  
router
- **network core:**
  - interconnected routers
  - network of networks



# The network edge

- end systems (hosts):
  - run application programs
  - e.g. Web, email
  - at “edge of network”
- client/server model
  - client host requests, receives service from always-on server
  - e.g. Web browser/server; email client/server
- peer-to-peer model:
  - minimal (or no) use of dedicated servers
  - e.g. VOIP, Skype



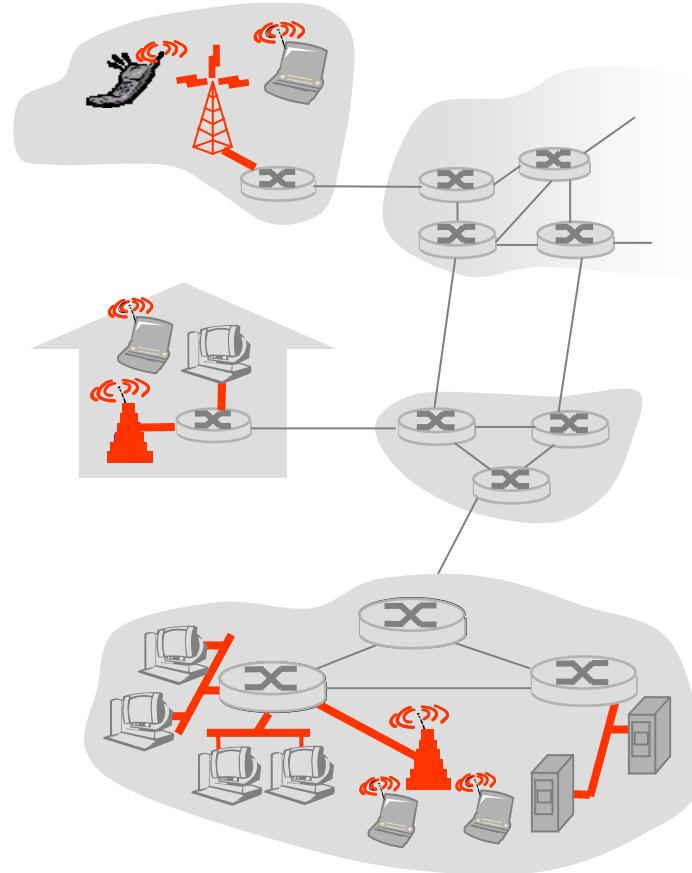
# Access networks

*How to connect end systems to edge router:*

- residential access nets
- institutional access networks (school, company)
- mobile access networks

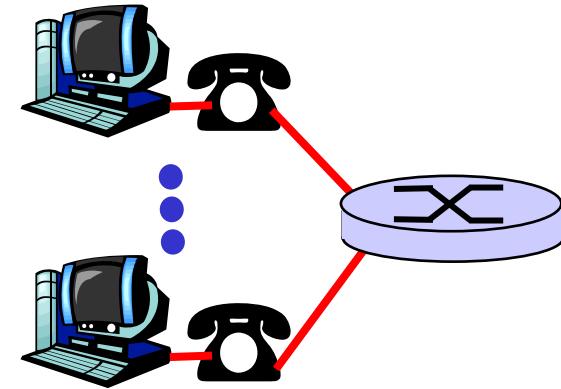
*Important parameters:*

- bandwidth (bits per second) of access network?
- shared or dedicated?



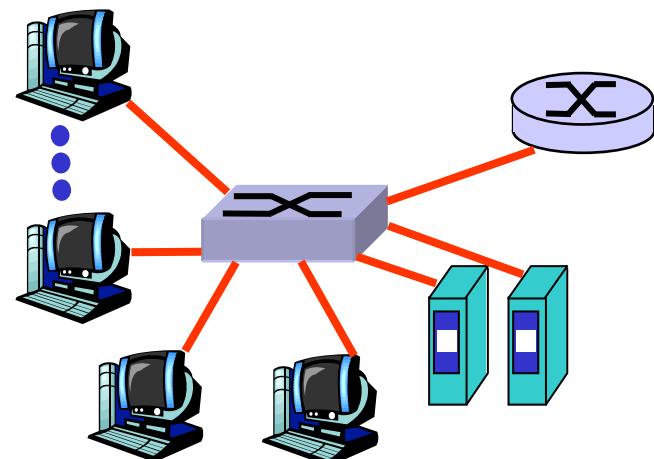
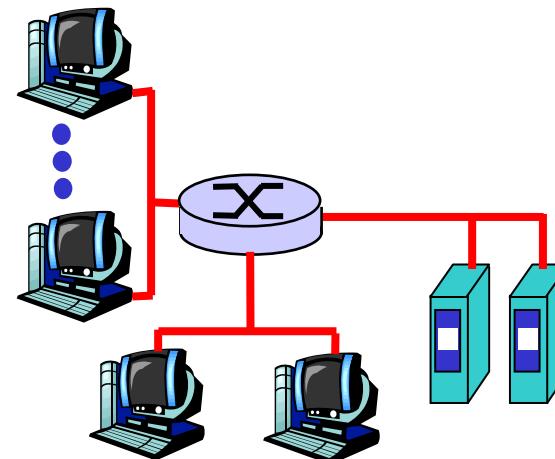
# Residential access: point-to-point access

- Dialup via modem
  - up to 56Kbps direct access to router (often less)
  - Can't surf and phone at same time: can't be “always on”
- DSL: digital subscriber line
  - deployment: telephone company or Internet provider
  - DSL divides communication link into 3 separate frequency bands (FDM):
    - up to 1 Mbps upstream (4 kHz bis 50 kHz)
    - up to 8 Mbps downstream (50 kHz bis 1 MHz)
    - normal telephone line (0 to 4 kHz)



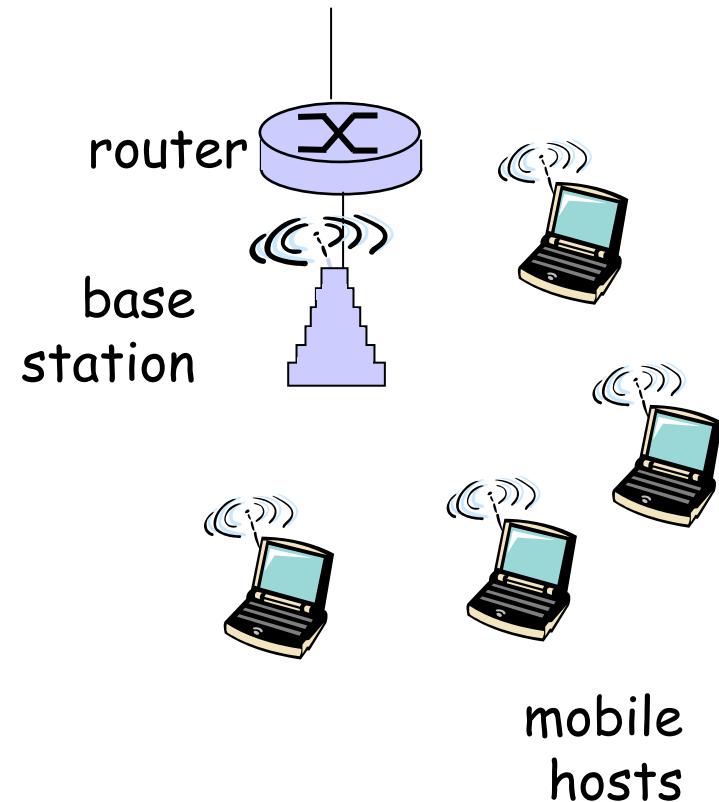
# Company access: local area networks

- company/univ **local area network** (LAN) connects end system to edge router
- **Ethernet:**
  - 10 Mbs, 100Mbps, 1Gbps, 10Gbps Ethernet
  - modern configuration: end systems connect into *Ethernet switch*



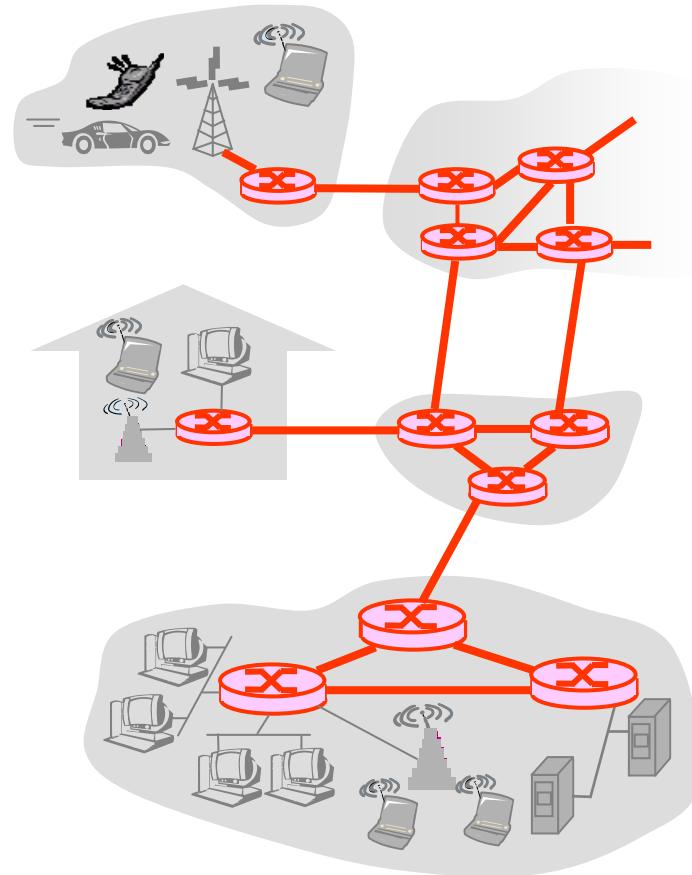
# Wireless access networks

- shared *wireless* access network connects end system to router
  - via base station (access point)
- wireless LANs:
  - 802.11b/g : 11 or 54 Mbps



# Network Core

- mesh of interconnected routers
- *the fundamental question:* how is data transferred through net?
  - **circuit switching:** dedicated circuit per call: telephone net
  - **packet-switching:** data sent thru net in discrete “chunks”



# Network Core: Circuit switching

End-to-end resources reserved for “call”

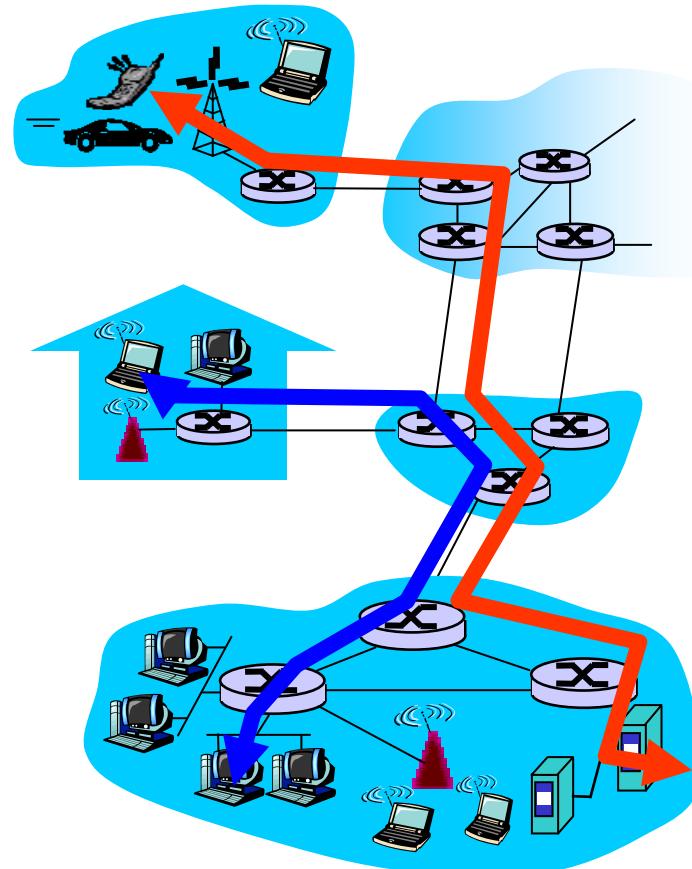
- link bandwidth, switch capacity
- dedicated resources: no sharing
- circuit-like (guaranteed) performance
- call setup required

network resources (e.g., bandwidth)  
divided into “pieces”

- pieces allocated to calls
- resource piece *idle* if not used by owning call (*no sharing*)

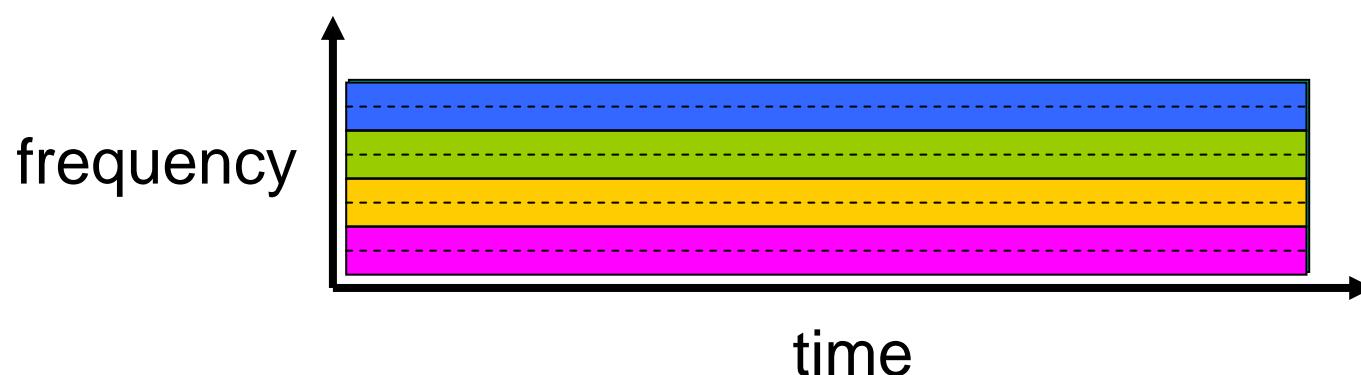
dividing link bandwidth into “pieces”

- frequency division
- time division

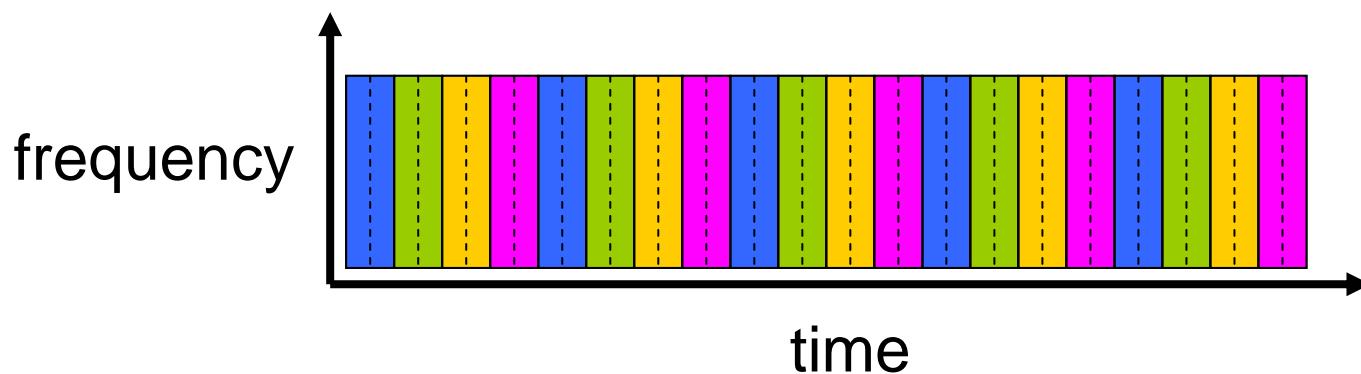


# FDM and TDM

FDM



TDM



Example:

4 users



# Network Core: Packet switching

each end-to-end data stream divided into  
*packets*

- user A, B packets *share* network resources
- each packet uses full link bandwidth
- resources used *as needed*

resource contention:

- aggregate resource demand can exceed amount available
- congestion: packets queue, wait for link use
- store and forward: packets move one hop at a time
  - switch receives complete packet before forwarding

Bandwidth division into ‘pieces’  
Dedicated allocation  
Resource reservation

