## Secrecy (1)

## Symmetric Key System:

Keys of Alice and Bob are identical and secret
Public Key System:
Both, Alice and Bob have a pair of keys, one is public, the other is only known by its holder.

## 1. Symmetric Key Systems (old)

Traditional encryption methods have been divided historically into two categories:

- substitution ciphers (preserve the order of the plaintext symbols but disguise them)
- transposition ciphers (reorders the plaintext symbols but do not disguise them)

Ancient and simple substitution cipher: Caesar's cipher
The ciphertext alphabet results from a shift of $k$ letters in the plaintext alphabet (key:=k).

Generalization of Caesar's chiffre: monoalphabetic substitution
Each letter or group of letters is replaced by another letter or group of letters to disguise it

## Example for a monoalphabetic substitution



## Secrecy (2)

## Transposition ciphers

Instead of disguising letters they are reordered

## Example for a columnar transposition


Plaintext pleasetransferonemilliondollarsto myswissbankaccountsixtwotwo
Ciphertext

## AFLLSKSOSELAWAIATOOSSCTCLNMOMANT

 ESILYNTWRNNTSOWDPAEDOBUOERIRICXB
## Symmetric Key Systems (1)

## 2. Symmetric Key Systems (modern)

Idea: Concatenation of standard transposition (permutation) and substitution elements (boxes):

Example for a P(ermutation)-box (01234567 ---> 36071245)

The order of sequence has changed


Example for a S(ubstitution)-box (3bit plaintext to 3bit ciphertext)
By appropriate wiring of the P-box inside, any substitution can be accomplished.
In this example:
Numbers $0,1,2,3,4,5,6,7$ each are replaced by the numbers 24506713


## Symmetric Key Systems (2)

Example for a product cipher (concatenation)


## Standard: DES

- plaintext is encrypted in blocks of 64 bits
- the algorithm has 19 steps
- the steps for decryption are done in the reverse order of those for encryption


## Public Key Systems (1)

## 3. Public-Key Systems

Basic problem behind:
Is it possible that Alice and Bob can communicate by encrypted messages without having exchanged before a common secret key?

Principal solution:
Each party has a pair of keys, a public one (accessible to everybody) and a private one (only known by itself)

## The general model



## Public Key Systems (2)

## The RSA algorithm

Two components:

- Selecting the keys
- Applying the encryption and decryption algorithm

Selecting the keys (by Bob):

1. Choose two large primes, $p$ and $q$
2. Compute $n=p \times q$ and $z=(p-1) \times(q-1)$.
3. Choose a number relatively prime to $z$, smaller than $n$ and call it $e$ ( $e$ is used for encryption).
4. Find $d$ such that $e \times d=1 \bmod z$ ( $d$ is used for decryption).

5 . The public key is ( $n, e$ ), the private key is ( $n, d$ ).

Encryption (by Alice) of a bit pattern (number) $m$ such that $m<n$ by means of Bob’s public key ( $n, e$ ).
The resulting cipher $c$ is:
$c=m^{e} \bmod n$

Decryption (by Bob) of $c$ by means of his private key $(n, d)$ in order to get the plaintext $m$ :

$$
m=c^{d} \bmod n
$$

## Public Key Systems (3)

## Example of the RSA algorithm

$p=5, q=7$---> $n=35, z=24$. Further, Bob selects $e=5, d=29$ (5*29-1 can be divided by 24)
----> public key of Bob: $(35,5)$, private key of Bob: $(35,29)$
Alice wants to send the message "LOVE" to Bob by encrypting each letter separately and interpreting each letter as the corresponding number ( a maps to $1, \ldots .$. , z maps to 26)

| Klartextbuchstabe | $m$ : numerische Darstellung | $m^{\text {e }}$ | Chiffretext $c=m^{e} \bmod n$ |
| :---: | :---: | :---: | :---: |
| L | 12 | 248832 | 17 |
| 0 | 15 | 759375 | 15 |
| V | 22 | 5153632 | 22 |
| E | 5 | 3125 | 10 |


| Chiffretext c | $c^{d}$ | Chiffretext $m=c^{d}$ $\bmod n$ | Kartext buchstabe |
| :---: | :---: | :---: | :---: |
| 17 | 481968572106750915091411825223072000 | 12 | 1 |
| 15 | 12783403948858939111232757568359400 | 15 | - |
| 22 | $8.51643319086537701195619449972111 e+38$ | 22 | v |
| 10 | 100000000000000000000000000000 | 5 | e |

## Authentication

## Authentication Protocols

- technique by which a process verifies that its actual communication partner is who it is supposed to be
- normally done before the partners start to exchange data messages, e.g. e-mails

Version with symmetric keys


## Version with public keys



